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# **EUROPEAN PATENT APPLICATION**

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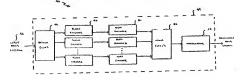
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#### (54)A method for high speed modulation and error control coding

(57)A digital data communication system is provided for transmitting a high speed differentially encoded data signal between a transmitter and a receiver. On the transmitter side (44), a high speed input data stream (42) is separated into a plurality of lower rate data streams by a demultiplexer (46). Prior to being multiplexed together to form a single high speed data stream, each of the lower rate data streams are individually block encoded by a group of Reed-Solomon block encoders (48) and then differentially encoded by a group of differential encoders (50). Alternatively, the plurality of low rate data streams are block encoded and then multiplexed together prior to being differentially encoded, so long as the multiplexing is performed on a code symbol-by-code symbol basis. Since the performance of a non-binary-based block code is based on an input code symbol error rate (as opposed to an input bit error rate), performance loss associated with either of the differential encoding schemes of the present invention are comparable to a non-differential encoding scheme.



# Description

# BACKGROUND OF THE INVENTION

# 1. Field of the Invention

[0001] This invention relates generally to a digital data communication system and, more particularly, to a differential encoding and error control coding scheme for high speed modulation in a digital data communication system.

# 10 2. Discussion of the Related Art

[10002] Modern data communication systems generally require low bit error probabilities. In a moderate-rate communication application (e.g., on the order of 100 Mb/s or less), a boward error correction scheme, such as comolutional encoding or block encoding, can be used to achieve low bit error rates (e.g., on the order of 10°) at low signal-to-noise ratios. However, a large number of mathematical operations per decoded bit are generally required for complex convotional codes (i.e., having high constraint length) or high overhead block codes (i.e., having a low ratio of information bits to parity-check high, As a result, the implementation of even moderate-rate decoding algorithms can be difficult in commercially viable hardware (e.g., CMOS).

[9003] In order to accommodate higher data rates (e.g., several hundred Mbks or more) while maintaining reasonable hardware complexity, multi-hop communication systems (where data is demodulated and remodulated at Intermediate locations prior to the final destination), generally go without coding and have reasonably simple modulation schemes. In some instances, vary light comodulational coding can be employed along with a greater transmit power and/or a larger antenna to overcome the degradations incurred when implementing a high speed demodulate-remodulate system. To achieve very low bit error probabilities (e.g. on the order of 10<sup>-0</sup>) at these higher data rates. It turns out at the Read-Solomon block codes are generally preferable to convolutional codes (of the same coding rate) because they provide better performance at the same signatio-noises ratio. Unfortunately, high rate (speed) decoding of a relatively modest complexity Read-Solomon block code is difficult (if not impossible) to implement in commercially available hardware.

[0004] Therefore, it is clearable to provide a digital data communication system for transmitting high speed data signals having very low bit error probabilities. By combining differential encoding with Read-Solomon block encoding and multiplexing techniques, degradations normally associated with differentially encoding are dramatically reduced. As a result, near-otherent performance can be achieved. Moreover, very low error probabilities can be obtained using low-complexity hardware, without the need to significantly after other system performance parameters, such as power consumption, size, and weight. The differential encoding scheme of the present invention also provides additional benefits to a coherently-detected high-rate system, such as resolving phase ambiguity and eliminating cycle-silogace problems.

# SUMMARY OF THE INVENTION

[9005] In accordance with the present invention, a digital data communication system is provided for transmitting as in pits period differentially encoded data signal between a transmitter and a receiver. On the transmitter side, a high speed input data stream its separated into a plurality of lower rate data stream by a demultiplexer. Prior to being multi-plexed back together to the ran a single high speed data stream, each of the lower rate data streams are individually continued by a group of Placed-Solomon (RS) block encoders and differentially encoded by a group of differential encoder. Atternatively, the plurality of low rate data streams are block encoded and them multiplexed together prior to long differential encoded, so long as the multiplexed receives an expectation of the prediction and prediction and prediction are comparable to a non-differentially encoded or here not detection scheme.

#### 50 BRIEF DESCRIPTION OF THE DRAWINGS

[0006] Other objects and advantages of the present invention will be apparent to those skilled in the art upon reading the following detailed description and upon reference to the drawings in which:

55 Figure 1 is a diagram illustrating a typical satellite data communication system in accordance with the present invention:

Figure 2 is a chart #fustrating how differential decoding typically causes two errors for each single error in the transmitted bit stream;

Figure 3 is a block diagram showing a demultiplexer operating on a bit-by-bit basis;

Figure 4 is a block diagram showing a demultiplexer operating on a code symbol-by-code symbol basis in accordance with a differential encoding scheme of the present invention;

Figure 5 illustrates the error pairs that can occur in an eight (8) bit block code symbol;

Figure 6 is a block diagram showing the demultiplexing operation being performed before the differential and block decoding operations in accordance with the present invention;

Figure 7 is a block diagram depicting a transmitter implementing a differential encoding scheme of the present invention;

Figure 8 is a block diagram depicting a receiver implementing the differential encoding scheme of the present invention; and

Figure 9 is a block diagram depicting a transmitter implementing an alternative differential encoding scheme of the present invention.

#### DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

[0007] Write the present invention is described herein with reference to illustrative embodiments for particular applications (e.g., inter-satellite inde), it should be undestood that the invention is not limited thereto. Any data communication system requiring very low bit error probabilities and high data rates without extreme bandwidth expansion could benefit from the present invention. Those having ordinary skill in the art and access to the leachings provided herein will recognize additional modifications, applications and embodiments within the scope thereof and additional fields in which the present invention would be of significant utility.

[9008] A typical satellite communication system 10 is depicted in Figure 1. The satellite communication system 10 includes at least one orbiting satellite 12 which may complete a virtual oftent connection between two of a plurality of ground stations 14. Generally, information is uplinked from a transmitting ground station to the satellite which in turn downlinks the information to a receiving ground station. However, in order to extend communication coverage of the satellite communication system 10, a satellite 12 may also communicate virtual rain inter-satellite link 16 directly with a second satellite 18 which may then communicate virtual rain statellite 18 which may then communicate virtual rains the satellite 18 which may then communicate virtual rains the satellite 18 which may then communicate virtual rains the satellite 18 which may then communicate virtual rains the satellite 18 which may then communicate virtual rains the satellite 18 which may then communicate virtual rains the satellite 18 which may then communicate virtual rains the satellite 18 which may the satellite 18 which may the communicate virtual rains the satellite 18 which may then communicate virtual rains the satellite 18 which may the satellite 18 which may then communicate virtual rains the satellite 18 which may be satellite 18 which may the satellite 18 which may be satellite 18 which may be satellite 18 which may the sate

[0009] For a typical high speed inter-satellite fink, the information rate is at least 500 Mb/s with a bit error rate requirement on the order of 10<sup>10</sup> or 10<sup>10</sup>. As will be more fully explained, a differential encoding scheme in accordance with the present invention is properly suited to handle the high speed nature of an inter-satellite ink requiring yet low bit error probability. In the context of a satellite communication system 10, the differential encoding scheme of the present invention is not exclusive to an inter-satellite link, it is envisioned that it may also be applicable to a high speed communication link between a single ground stations and a satellite, or between two terrestrial ground stations.

[0010] With regards to bit error rate performance, differential encoding pombined with coherent demodulation (or detection) generally yields approximately twice as many errors as coherent detection alone. When used to make a differential decision, a single error in a transmitted bit stream causes two errors in the differentially decoded bit stream for the bits in gray denote errors is a shown in Figure 2. More specifically, the bit error probability of a non-differentially encoded-coherently detected binary phase-shift keying (BFS) data signal is given by:

$$P_b(E) = O\left(\sqrt{\frac{2E_b}{N_o}}\right)$$

45 whereas, the bit error probability of a differentially encoded-coherently detected BPSK signal is given by:

$$P_b(E) = 2O\left(\sqrt{\frac{2E_b}{N_o}}\right)\left[1-O\left(\sqrt{\frac{2E_b}{N_o}}\right)\right]$$

Trus, the error probability for differential encoding-coherent detection is approximately twice the error probability for non-differentially coder-coherent detection. However, some of the performance loss can be regional depending on how and where the differential encoding is performed in relation to the other signal processing techniques (e.g., block-encoding and multiplexing) employed in the statellise link. Accordingly, the differential encoding scheme of the prescription achieves bit error rates comparable to those associated with non-differential encoding-coherently detected schemes.

[0011] In Figure 3, a demultiplexer (1:3) 20 receives an input data signal 16 and seperates it into three output sig-

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nals 18. When the demultiplexer 20 operates on a bit-by-bit basis, each bit going into the demultiplexer 20 is placed on the next subsequent output line. For a non-differentialty encoded data stream, the bit error rate and code symbol error rate going into the demultiplexer 20 will be identical to the bit and code symbol error rates going into each of a series of three block decoders 22. In the case of an input data signal containing fifteen (15) errors for every 3000 bits of data, each output fine exhibits as bit error rate of \$5.70 in.

[0012] In contrast, if the input signal 16 going into the demultiplexer 20 is a differentially encoded data signal, then it will contain approximately 90 ornors for every 2000 bits of data. Although these 30 errors generally occur in pairs the pairs are randomly distributed over AB, AC, or BC bit combinations, such that the net effect is to doubt in amount of errors going into each of the block decoders 22. Assuming that the error pairs are generally spaced more than 8 bits apart (for a block code based on 8 bit symbols), the code symbol error rate going into each of the block decoders 22 is also doubted.

[0013] Next, consider the case where a demultiplower 24 operates on a code symbol-by-code symbol basis. In Figure 4, the demultiplemer (1) 32 or receives a differentially decoded input data signal 16 and seperates it into three outsisionals 18. For illustration purposes, the code symbol is based on an 8 bit or byte symbol glato referred to herein as a byte-by-byte basis). Thus, each sequence of 8 bits (i.e., a code symbol) going into the demultiplemer 24 separates on the next subsequent cultural line. Because the demultiplemer 24 operates on code symbol basis, it is likely that each of the error pairs will be crounced within a block code symbol.

[0014] Figure 5 illustrates how there is a 7 in 9 chance that an error pair will end up in the same block whole, in the event that the error pair falls within the same symbol, the symbol error rate of the falls oping in block of the block decoders 22 does not increase. On the other hand, 2 out of 9 times the error pair falls on a symbol boundary, thereby increasing the symbol error rate going into each of the block decoders 22. One stelled in the art will readily recognize from such discussions, that this differential encoding scheme is applicable to any block code having at least 2 bits per code symbol.

[0015] Since the performance of a block code is based on its input code symbol error rate and not its input bit error rate (although the two are related, it is not one-house and unique, the performance loss associated with demultipling a differentially decoded data signal on a symbol by-symbol basis (i.e. block code symbol), is minimal when compared to a non-differentially encoded data signal. The following table summarizes the performance loss values:

	Encoding Method	Demux Method	Req'd E <sub>b</sub> /N <sub>o</sub> (dB)	Loss (dB)			
Bits	Non-differential (coherent)	bis-by-bit	5.692	0 (reference)			
	Non-differential (coherent)	byte-by-byte	6.692	0			
	Differential	bit-by-bit	7.324	0.632			
	Differential	byte-by-byte	6.867	0.175			

[0016] Overall, this type of differential encoding scheme exhibits some degradation (0.175 dB) in relation to strictly non-differential encoding (0 dB), and considerably less degradation than demultiplexing a differentially encoded data signal (0.532 dB) on a bit-by-bit basis.

[0017] It should also be noted that bit error rate performance is stightly different when the above-described differential encoding scheme is applied to quadrather phase-shit keying (QPSi) modulated signals. As will be appared to one skilled in the art, the multipleare operates on a QPSiK symbol-by QPSiK symbol basis, but otherwise the error probability for QPSiK symbol-by-QPSiK symbol demultiplearing is anatigous to the above-described case of bit-by-bit demultiplearing. However, when four QPSiK symbols are demultipleared together (analogous to byte-by-byte demultiplearing), there is a 3 in 5 chance that two adjacent QPSiK symbols in error will lie within the same block symbol, and therefore not increase the error probability of the data going into the block decoders. Accordingly, 2 out of 5 times the errors will lie on the block symbol boundary, and thus increase the error probability of the data going into the block symbol boundary, and thus increase the error probability of the data going into the block symbol boundary, and thus increase the error probability of the data going into the block symbol collowing table summarizes the performance loss values:

55		Encoding Method	Demux Method	Reg'd E <sub>b</sub> /,N <sub>e</sub> (dB)	Loss (dB)		
	QPSK Symbols	Non-differential (coherent)	symOpsk-by-symOpsk	6.692	O (reference)		

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#### (continued)

Encoding Method	Demax Method	Reg'd E <sub>b</sub> /N <sub>o</sub> (dB)	Loss (dB)			
Non-differential (coherent)	sm <sub>RS</sub> -by-sym <sub>RS</sub>	6.692	0 -			
Differential	sym <sub>QPSK</sub> -by-sym <sub>QPSK</sub>	7.324	0.632			
 Offerential	sym <sub>RS</sub> -by-sym <sub>RS</sub>	6.987	0.295			

[0018] Again, there is considerably less degradation associated with demultiplexing a differentially decoded data signal on a code symbol-by-code symbol basis (0.295 dB) as compared to demultiplexing on a OPSK symbol-by-

[0019] In each of the above cases, differential encoding is performed after multiplexing on the transmitter side. However, similar performance loss can be achieved by moving a series of differential encoders in front of the multiplexer. Accordingly, a series of differential decoders 30 is moved after a demultiplexer 32 on the receiver side as shown in Figure 6. First, the demultiplexer (1.3) 32 receives an input data signal 34 from the transmitter and separates it into three output signals 36. Next, each of the output signals are differentially decoded by the series of three lockle decoders 38.

[0020] In this alternative differential encoding scheme, the error pairs that arise from differential decoding with adveys behave in the same manner as the proviously described by the-ty-tyet case. That is, there will be a 7 in 0 chance that the error pair will fall within the same block symbol, and only 2 out of 9 times will the error pair fall on a symbol boundary. As a result, multiplexing on a bit-ty-but for a symbol-by-symbol basis has no effect on performance loss to long as differential encoding is performed before multiplexing on the transmitter side. The performance loss for this alternative differential encoding scheme is on the order of 0.1758. Moving the differential encoding scheme is on the order of 0.0758. Moving the differential encoding is not the order of 0.0758 which is the performance loss is on the order of 0.0758 of 0.0758 which is the performance loss is on the order of 0.0758 which is the

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10021] A first preferred embodiment of a differential encoding scheme 40 for use on the transmitter side of a high speed satellite link is shown in Figure 7. Due to the complexity of implementing a single 500 Mb/s block decopiting an imput data stream 42 into a transmitter 44 is split up into N lower rate data stream by a (1:1) detuntiplexer 46.05ptt ing up the encoding and decoding process into several parallel data streams reduces the hardware complexity of the decoders which carn in turn operate at a lower rate without any loss in performance over a single, high speed decoder. As will be apparent to one skilled in the art, the implementation of the demultiplexer may vary in accordance with the present investor.

[0022] Next, each of the lower rate data streams undergoes block encoding by a set of block encoders 48, auch that seach of the block encoders 48 operates at 11N times the speed of the overall data rate of the satellite link. As previously discussed, the use of convolutional codes is precluded by the complexity of high speed decoders as well as by the bandwidth expansion necessary to active every low bit enror rates. As a result, a block code was chosen for enough each of the lower rate data streams. More specifically, the byte organized ATM cell structure associated with a satelline inlik lends tissed to a (255,239) Reed-Solomon code having 95 the (i.e., byte size) symbols. Although a Reed-Solomon block code is preferable, it is envisioned that other types of block codes (based on symbols having at least 2 bits per symbol) may also be implemented in accordance with the present invention.

[9023] Each of the separately block encoded data streams are then differentially encoded by a set of differential encoders 60, and then multiplexed together by a (N:1) multiplexed 52 to form a single high rate transmissible data stream 56. As will be apparent to one skilled in the art, this data stream is preferably OPSK modulated by a modulator 54 before being transmitted by the transmitter. It is also envisioned that other modulation techniques may be applied to the data stream prior to its transmission by the transmitter.

[9024] On the receiver side 60, the transmitted data stream 56 is likewise demultiplexed, differentially decoded and then block decoded as shown in Figure 8. First, the transmitted data stream 56 is demodulated by a demodulater 62. Next, the demodulater 62 is expented into a plurality of demodulated data streams by a Gemultiplexer (1:N) 64. Each of the demodulated data streams by a Gemultiplexer (1:N) 64. Each of the demodulated data streams are then differentially decoded by a series of differential decoders 68 is demodulated data stream.

[6025] In an alternative embodiment, each of the block encoded data streams are multiplexed together and then of differentially encoded as shown in Figure 9. More specifically, on the transmitter side 70, an input data stream 72 is separated in a plurality of data streams by a demultiplexer 74. Each of the data streams are the block encoded by a series of block encoders 76 prior to being combined into a single data stream by a multiplexer 78. Next. a differential encoder 80 differentially encodes the data stream. Lastly, the data stream is modulated by a modulator 82. Since other

modulators naturally encode data in a differential manner, this embodiment is typically preferred in a optical-based (rather than RP) link. Or the receiver side (not shown), demulgistering is performed on a code symbol-by-code symbol basis after differential decoding. Otherwise, this afternative embodiment of the differential encoding scheme of the present invention is similar to the lists preferred embodiment.

5 [0065] It should be appreciated that the differential encoding scheme of the present invention is properly suited to handle the high speed nature of an inter-satellite fink requiring very low bit error probabilities can be obtained using low-complexity, commercially available hardware without the need to significantly after other performance parameters, such as power consumption, size, and weight. Furthermore, the differential encoding scheme of the present invention others several additional benefits, such as removing the phase armbiguity that arises from 1p hase-shift-keyed signaling (e.g., a phase ambiguity of pl/2 for QPSK) and alleviating cycle-slipping that could otherwise degrade the performance of a non-differentially encoded system.

[0027] The foreign discloses and describes merely exemplary embodiments of the present invention. One skilled in the str will readily recognize from such discussion, and from the accompanying drawings and claims, that value changes, modifications and variations can be made therein without departing from the spirit and scope of the present invention.

#### Claims

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- A digital data communication system for transmitting a high speed differentially encoded data signal between a transmitter and a receiver, the transmitter comprising;
  - a plurality of block encoders each receiving one of a plurality of data signals and applying a block code to form one of a plurality of block coded data signals:
  - a plurality of differential encoders each receiving one of said plurality of block coded data signals and applying a differential code to form one of a plurality of differentially coded data signals; and
  - a multiplexer receiving said plurality of differensatily coded data signals and combining said plurality of differentially coded data signals to form a differentially coded transmissible data signal.
  - The digital data communication system of Claim 1 wherein said plurality of block encoders applies a block code based on symbols having at least two bits.
  - The digital data communication system of Claim 1 wherein said plurality of block encoders applies a Reed-Solomon block code.
- 35 4. The digital data communication system of Claim 1 further comprising a demultiplexer receiving an input data signal and securating said input data signal into said plurality of data signals.
  - The digital data communication system of Claim 1 further comprising a modulator receiving said differentially coded transmissible data signal and modulating said differentially coded transmissible data signal prior to transmission by the transmitter.
    - The digital data communication system of Claim 1 wherein the receiver of said differentially coded transmissible data signal includes:
- 45 a demodulator receiving said differentially coded transmissible data signal for demodulating said differentially coded transmissible data signal into a demodulated data signal:
  - a demultiplexer receiving said demodulated data signal for separating said demodulated data signal into a plurality of demodulated data signals;
  - a plurality of differential decoders each receiving one of said plurality of demodulated data signals for differentially decoding into a plurality of differentially decoded data signals;
    - a plurality of block decoders each receiving one of said plurality of differentially decoded data signals for decoding into a plurality of data signals; and
    - a second multiplexer receiving said plurality of data signals and combining said plurality of data signals to form an output data signal.
  - 7. The digital data communication system of Claim 1 wherein the digital data communication system being a satellite communication system, such that the transmitter being associated with at least one of a satellite and an earth terminal and the receiver being associated with at least one of a second satellite and a second earth terminal.

- A digital data communication system for transmitting a high speed differentially encoded data signal between a transmitter and a receiver, the transmitter comprising:
  - a plurality of block encoders each receiving one of a plurality of data signals and applying a block code to form one of a plurality of block coded data signals, said block code being based on symbols having at least two bits; a multiplexer receiving said plurality of block coded data signals and combining said plurality of block code data signals to form a multiplexed data signal, said multiplexer operating on a block code symbol-by-block code symbol basis; and
  - a differential encoder receiving said multiplexed data signal and applying a differential code to form a differentially coded data signal.
- A method for transmitting a high speed differentially encoded data signal between a transmitter and a receiver in a digital data communication system, comprising the steps of:
  - applying a block code to each of a plurality of data signals to form a plurality of block coded data signals; applying a differential code to each of said plurality of block coded data signals to form a plurality of differentially coded data signals; and
  - multiplexing each of said plurality of differentially coded data signals to form a differentially coded transmissible data signal.
- 10. A method for transmitting a high speed differentially encoded data signal in a digital data communication system, comprising the steps of:
  - applying a block code to each of a plurality of data signals to form a plurality of block coded data signals, said block code being based on symbols having at least two bits;
  - multiplesting each of said plurality of block coded data signals on a block code symbol-by-block code symbolby-block code symbol beach; thereby forming a multiplexed data signal; and applying a differential code to said multiplexed data signal to form a differentially coded transmissible data signal
- nal.
- 11. The method of Claim 10 further comprising the steps of:
  - receiving said differentially coded transmissible data signal at the receiver:
- demodulating said differentially coded transmissible data signal into a demodulated data signal; differential decoding said demodulated data signal to form a differentially decoded data signal;
  - demultiplexing said differentially decoded data signal into a plurality of differentially decoded data signals on a block code symbol-by-block code symbol basis:
    - applying a block decoding to each of said differentially decoded data signals to form a plurality of decoded data signals; and
- multiplexing each said plurality of decoded data signals into an output decoded data signal.

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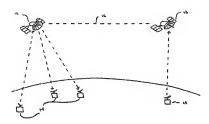
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FIGURE Z

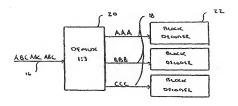


FIGURE 3

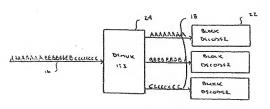


FIGURE 4

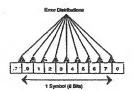


FIGURE S

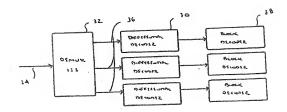
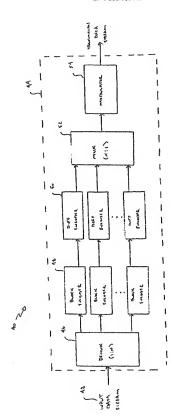
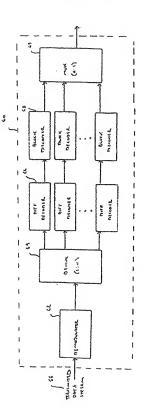


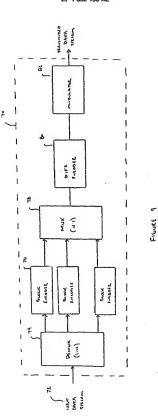
FIGURE 6



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